# CS316: Compilers Lab

Programming Assignment 2: Parser, Due: 31/1/2023

### 1 Introduction

Your goal in this step is to generate the parser for your programming language's grammar. By the end of this step your compiler should be able to take a source file as the input and parse the content of that file returning "Accepted" if the source file's content is correct according to the grammar or "Not Accepted" if it is not.

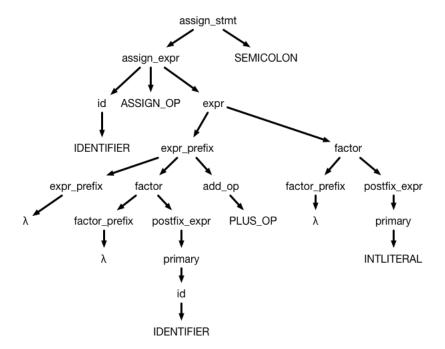
Now the scanner created earlier will be modified to feed the parser. Instead of printing the tokens, the scanner has to return what token is recognized in each step.

## 2 Background

The job of a parser is to convert a stream of tokens (as identified by the scanner) into a parse tree: a representation of the structure of the program. So, for example, a parser will convert:

$$A := B + 4$$
;

into a tree that looks something like:



This tree may look confusing, but it fundamentally captures the structure of an assignment statement: an assignment expression followed by a semicolon. An assignment expression is decomposed into an identifier followed by an assignment operation followed by an expression. That expression is decomposed into a bunch of primary terms that are combined with addition and subtraction, and those primary terms are decomposed into a bunch of factors that are combined with multiplication and division (this weird decomposition of expressions captures the necessary order of operations). Eventually, those factors, through a series of derivations, become identifiers or literals.

One important thing to note is that the *leaves* of the tree are the tokens of the program. If you read the leaves of the tree left to right (ignoring lambdas, since lambdas just represent the empty string), you get:

#### IDENTIFIER ASSIGN\_OP IDENTIFIER PLUS\_OP INTLITERAL SEMICOLON

Which is exactly the tokenization of the input program!

### 2.1 Context-free grammars

To figure out how each construct in a program (an expression, an if statement, etc.) is decomposed into smaller pieces and, ultimately, tokens, we use a set of rules called a *context-free grammar*. These rules tell us how constructs (which we call "non-terminals") can be decomposed and written in terms of other constructs and tokens (which we call "terminals").

#### 2.2 Micro

The context-free grammar that defines the structure of Micro (the name of your programming language) is:

```
/* Program */
program
                  -> PROGRAM id BEGIN pgm body END
id
                  -> IDENTIFIER
                  -> decl func_declarations
pgm body
                    -> string decl decl | var decl decl | empty
decl
/* Global String Declaration */
string\_decl
                 \rightarrow STRING id := str ;
                  -> STRINGLITERAL
str
/* Variable Declaration */
var decl
                  -> var type id list ;
var\_type
                   -> FLOAT | INT
any\_type
                  -> var_type | VOID
id_list
                  -> id id tail
id tail
                  -> , id id tail | empty
/* Function Parameter List */
param decl list -> param decl param decl tail | empty
param decl
                  -> var type id
param decl tail -> , param decl param decl tail | empty
/* Function Declarations */
func declarations -> func decl func declarations | empty
                 -> FUNCTION any type id (param_decl_list) BEGIN func_body END
func decl
func\_body
                  -> decl stmt list
/* Statement List */
stmt list
                  -> stmt stmt list | empty
\operatorname{stmt}
                  -> base stmt | if stmt | while stmt
base stmt
                  -> assign stmt | read stmt | write stmt | return stmt
/* Basic Statements */
assign stmt
                 -> assign expr ;
                 \rightarrow id := expr
assign expr
                 -> READ ( id list );
read stmt
                 -> WRITE ( id list );
write stmt
return stmt
                 -> RETURN expr ;
```

```
/* Expressions */
                  -> expr prefix factor
expr
\exp r\_prefix
                  -> expr prefix factor addop | empty
factor
                  -> factor_prefix postfix_expr
factor_prefix
                  -> factor prefix postfix expr mulop | empty
postfix expr
                  -> primary | call expr
call expr
                  -> id ( expr list )
{\tt expr\_list}
                  -> expr expr list tail | empty
expr list tail
                  -> , expr expr list tail | empty
                  -> ( expr ) | id | INTLITERAL | FLOATLITERAL
primary
                  ->+ | -
addop
mulop
                  _> * | /
/* Complex Statements and Condition */
                  -> IF ( cond ) decl stmt list else part ENDIF
if stmt
else\_part
                  -> ELSE decl stmt list | empty
cond
                  -> expr compop expr
                  -> < | > | = | != | <= | >=
compop
                  -> WHILE ( cond ) decl aug stmt list ENDWHILE
while stmt
aug stmt list
                  -> aug stmt aug stmt list | empty
                  -> base stmt | aug if stmt | while stmt | CONTINUE; | BREAK;
aug stmt
                  -> IF ( cond ) decl aug stmt list aug else part ENDIF
aug if stmt
aug else part
                  -> ELSE decl aug stmt list | empty
```

So this grammar tells us, for example, that an if\_stmt looks like the keyword IF followed by an open parenthesis, followed by a cond expression followed by some decl (declarations) followed by a stmt\_list followed by an else\_part followed by the keyword ENDIF.

An input program matches the grammar (we say "is accepted by" the grammar) if you can use the rules of the grammar (starting from program) to generate the set of tokens that are in the input file. If there is no way to use the rules to generate the input file, then the program does not match the grammar, and hence is not a syntactically valid program.

## 3 Building a Parser

There are many tools that make it relatively easy to build a parser for a context free grammar (in class, we will talk about how these tools work): all you need to do is provide the context-free grammar and some actions to take when various constructs are recognized. The tools we recommend using are:

- Bison (this is a tool that is meant to work with scanners built using Flex). Note that integrating a flex scanner with bison requires a little bit of work. The process works in several steps that seem interlocking:
  - Define your token names as well as your grammar in your bison input file (called something like microParser.y)
  - Run bison -d -o microParser.cpp on microParser.ypp, it will create two output files: microParser.tab.cpp (which is your parser) and microParser.tab.hpp (which is a header file that defines the token names)
  - 3. In your scanner file (called something like microLexer.1), add actions to each of your token regexes to simply return the token name (you define the token names in your .y file). (Warning: make sure that you don't have a token named BEGIN even if the regex matches the string "BEGIN", because that will cause weird, hard-to-find errors. Call that token something like \_BEGIN.) To

make sure that your scanner compiles, you will need to put #include 'microParser.tab.hpp' in the part of your .1 file where you can include C code.

- 4. Run Flex on microLexer.1 to produce lex.yy.cpp
- 5. In another file, main.c/cpp, write a main function (this file will also need to include microParser.tab.hpp). Your main function should open the input file and store the file handle in a variable called yyin. Calling yyparse() will then run your parser on the file associated with yyin.
- Compile together all of microParser.cpp, lex.yy.cpp, and your main.c/cpp function to build your compiler.
- ANTLR (this is the same tool that can also build lexers). You should define your grammar in the same .g4 file in which you defined your lexer.
  - 1. Running ANTLR on that .g4 file will produce both a Lexer class and a Parser class.
  - 2. In your main file, rather than initializing a lexer and then grabbing tokens from it (as you may have done in step 1), you instead initialize a lexer, initialize a CommonTokenStream from that lexer, then initialize a parser with the CommonTokenStream you just created.
  - 3. You can then call a function with the same name as your top-level construct (probably program) on that parser to parse your input.

Note that the names of the files used here are for example purposes only. You may choose any filename.

### 4 What you need to do

You must build upon your scanner's implementation done for PA1.

The grammar for Micro is given above. All you need to do is have your parser parse the given input file and print Accepted if the input file correctly matches the grammar, and Not Accepted if it doesn't (i.e., the input file cannot be produced using the grammar rules).

In Bison's parser, you can define a function called yyerror (look at the documentation for the appropriate signature) that is called if the parser encounters an error.

In ANTLR, this is a little more complicated. You will need to create a new "error strategy" class (extend DefaultErrorStrategy) that overrides the function reportError. You can then set this as the error handler for your parser by calling setErrorHandler on your parser before starting to parse.

Sample inputs and outputs are provided to you along with the starter files that you need to get started on this assignment.

## 5 What you need to submit

- Place all the necessary code for your compiler (including the .1 file) that you wrote *yourself*. You do not need to include the ANTLR jar files if you are using ANTLR.
- A Makefile with the following targets:
  - 1. compiler: this target will build your compiler
  - 2. clean: this target will remove any intermediate files that were created to build the compiler
  - 3. dev: this target will print the same information that you printed in previous PA.
- A shell script (this must be written in bash) called runne that runs your compiler (scanner+parser). This script should take in two arguments: first, the input program file to be compiled and second, the filename where you want to put the compiler's/parser's output. You can assume that we will have run make compiler before running this script.

• You should tag your programming assignment submission as cs316pa2submission

While you may create as many other directories as you would like to organize your code or any intermediate products of the compilation process, both your Makefile and your runme script should be in the root directory of your repository.

Do not submit any binaries and do not delete/rename any of the files provided to you as part of the starter set. Your git repo should only contain source files; no products of compilation.